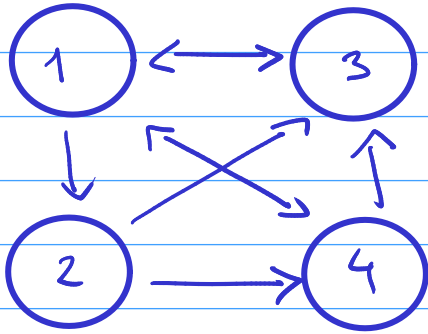


PageRank algorithm

Motivation



Graph: 4 nodes
adjacency matrix

$$\text{Adj} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix}$$

Q. which node is the most important?

- "outdegree": ① has 3
- "indegree" / backlinks: ③ has 3

Markov chain

before the transition matrix:
$$P_{ij} = \frac{\text{Adj}(i,j)}{\sum_{h=1}^n \text{Adj}(i,h)}$$

↳ row normalized

ex:

$$P = \begin{bmatrix} 0 & 1/3 & 1/3 & 1/3 \\ 0 & 0 & 1/2 & 1/2 \\ 1 & 0 & 0 & 0 \\ 1/2 & 0 & 1/2 & 0 \end{bmatrix}$$

P is a stochastic matrix:

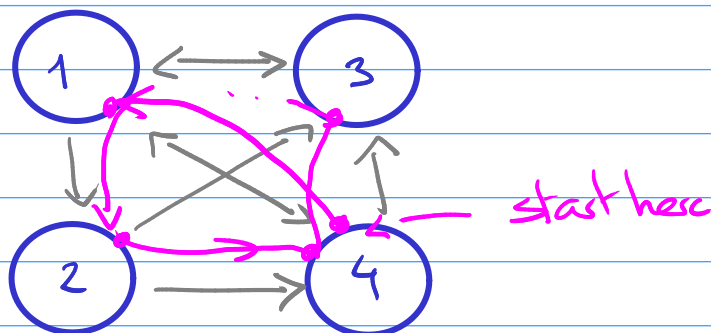
- $P_{ij} \geq 0$
- $\sum_{j=1}^n P_{ij} = 1$ for all i

We can define a Markov chain: X_n "jumping" on the graph

with $P(X_{n+1} = j | X_n = i) = p_{ij}$

↳ x jumping from node i to node j

ex: $x_1 = 4, x_2 = 1, x_3 = 2, x_4 = 4, x_5 = 3, \dots$



Remark: The paper uses the notation:

$$A = P^T \quad (\text{Transpose of } P)$$

↳ A is column-stochastic matrix.

Idea: The most important node is the one that "appears" the most.

↳ run code

Q: is there a fast way to find the "proportion"?

Equilibrium distribution

Denote q_n the probability density vector of X_n :

$$q_n(i) = \mathbb{P}(X_n = i)$$

ex: if $X_1 = 4$ always, Then $q_1 = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}$

Then X_2 is 1 or 3 with probability $\frac{1}{2}$, Thus:

$$q_2 = \begin{bmatrix} 1/2 \\ 0 \\ 1/2 \\ 0 \end{bmatrix}$$

How to get from q_1 to q_2 ? $q_2 = P^T q_1$.

More generally:

$$q_{n+1} = P^T q_n$$

$\leadsto$

extra
"forward Kolmogorov eq."

$$\partial_t p = L^* p$$

"backward Kolmogorov eq"

$$\partial_t u = L u$$

If the dynamics converges $q_n \rightarrow q_\infty$,

Then we have a fixed point:

$$q_\infty = P^T q_\infty$$

$\Rightarrow q_\infty$ eigenvector of $P^T (=A)$

associated with $\lambda = 1$.

$\hookrightarrow$ run code

ex: $q_\infty = \begin{bmatrix} .39 \\ .13 \\ .29 \\ .19 \end{bmatrix}$ $\leftarrow$

Linear algebra

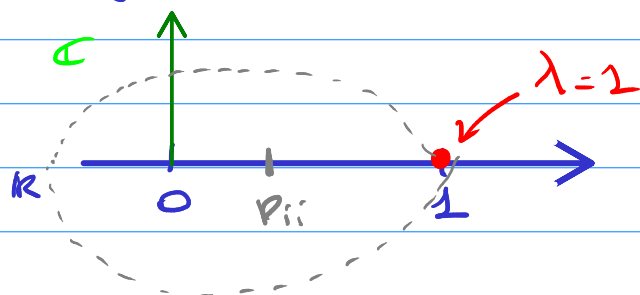
Questions: 1. Is ρ_∞ unique? i.e. $\lambda=1$ has multiplicity 1?

←

2. How to compute it for very large matrix?

Perron
Frobenius

spectrum of P (or P^T)



Since P stochastic matrix

$$P \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix} = 1 \cdot \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

Gershgorin's Thm:

$$\lambda \in \bigcup_i B(p_{ii}, \sum_{j \neq i} |p_{ij}|)$$

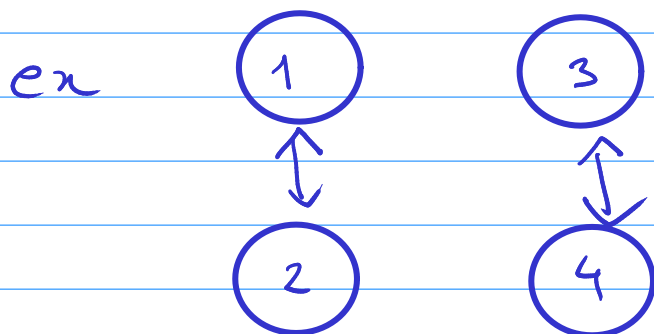
$$P = \begin{bmatrix} \dots & p_{ii} & \dots \end{bmatrix}$$

$$\text{Then: } \operatorname{Re}(\lambda) \leq 1.$$

$\Rightarrow \lambda=1$ is the "largest" eigenvalue.

Problem: it can have multiplicity > 1 if the graph

is "disconnected" (i.e. P "reducible")



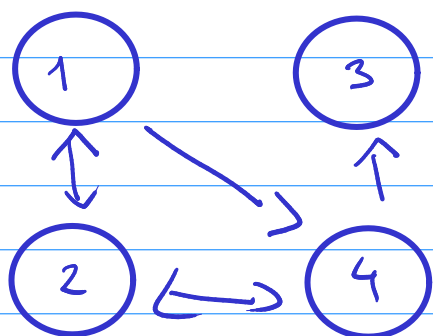
$$\text{Adj} = \begin{bmatrix} \boxed{0 & 1} & 0 & 0 \\ \boxed{1 & 0} & 0 & 0 \\ 0 & 0 & \boxed{0 & 1} \\ 0 & 0 & \boxed{1 & 0} \end{bmatrix}$$

Two eigenvectors associated with $\lambda=1$:

$$u_1 = \begin{pmatrix} 1/2 \\ 1/2 \\ 0 \end{pmatrix} \quad \text{and} \quad u_2 = \begin{pmatrix} 0 \\ 0 \\ 1/2 \\ 1/2 \end{pmatrix}.$$

Another problem: "absorbing states"

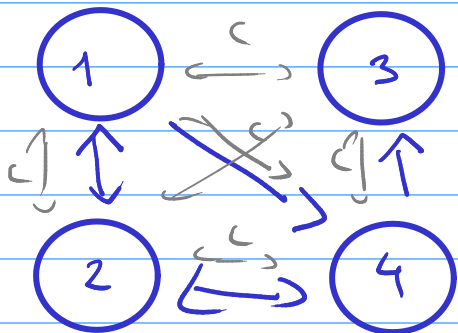
$\hookrightarrow$ dangling node



$\leftarrow$ cannot escape (3)

Thus: $q_\infty = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}.$

To avoid these 2 problems, add " ϵ " connexion everywhere



$$P^\epsilon = (1-\epsilon)P + \epsilon \begin{bmatrix} 1/n & \dots & 1/n \\ 1/n & \dots & 1/n \end{bmatrix}$$

$\hookrightarrow$
 stochastic
 + "irreducible"
 + a-periodic

} Perron Frobenius $\exists! q_\infty$ s.t.

$$q_\infty = P^{\epsilon T} q_0$$

with $q_0 > 0$ (and $|q_0| = 2$)

How to compute q_∞ ?

$$q_{n+1} = \mathbb{E} q_n \quad \text{vs} \quad q_n \xrightarrow{n \rightarrow \infty} q_\infty.$$

$\uparrow$
exponentially
(given by the spectral gap

$$" \lambda_1 - \lambda_2 "$$

"1"

$\hookrightarrow$ link with Fiedler number

$\hookrightarrow$ Algebraic connectivity

$\hookrightarrow$ run code.